

23/09/2021

Lecture 2

$\vec{x} \rightarrow$ column vector in $\mathbb{R}^n$

Linear Programs:

$$\begin{array}{l} \text{minimize} \quad f_0(\vec{x}) = \vec{c}^T \vec{x} \\ \text{subject to} \quad f_i(\vec{x}) \leq b_i, \quad i=1, \dots, m \\ f_i\text{'s are linear maps from } \mathbb{R}^n \text{ to } \mathbb{R}. \\ \quad \quad \quad \hookrightarrow \text{constraint functions} \end{array}$$

More concretely,

$(x_1), \dots, (x_n) \rightarrow$ decision variables

minimize $c_1 x_1 + \dots + c_n x_n \quad (= \vec{c}^T \vec{x})$

subject to

$$a_{11} x_1 + \dots + a_{1n} x_n \leq b_1$$

$$\bar{a}_{21} x_1 + \dots + \bar{a}_{2n} x_n \leq b_2$$

$$\vdots$$

$$a_{m1} x_1 + \dots + a_{mn} x_n \leq b_m$$

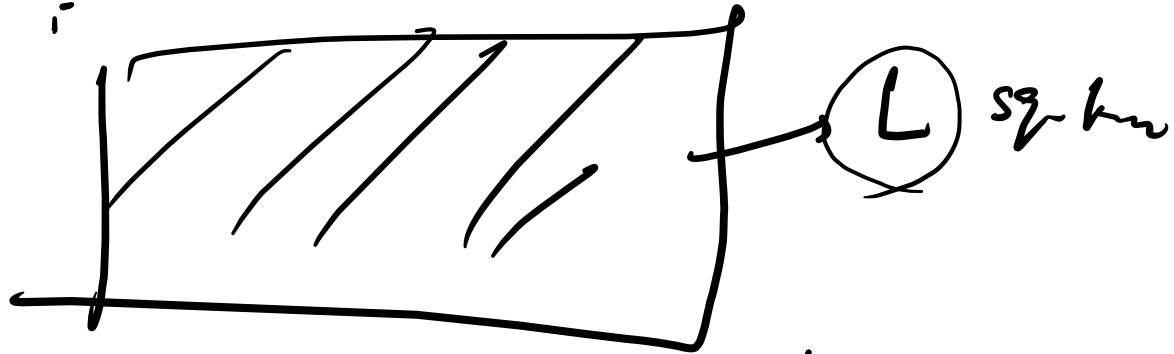
$$x_1 + x_2 = 1$$

$$x_1 + x_2 \leq 1$$

$$(-x_1) + (-x_2) \leq -1$$

Examples of LP ;

Example 1 :



Amount of fertilizer $\rightarrow \underline{F}$ kg.
" " pesticide $\rightarrow P$ kg.

Fertilizer req.
(per sq. km.)

Pesticide req.
(per sq. km.)

selling price
(per sq. km.)

Wheat

F_W

P_W

$\underline{\underline{S_W}}$

Mary

F_M

P_M

$\underline{\underline{S_M}}$

Area of land allocated to W $\rightarrow x_W$
" " M $\rightarrow x_M$

"

maximize: $\boxed{S_W x_W + S_M x_M} \leq -\alpha$

subject to: $\underline{x_W + x_M} \leq \underline{L}$ (total land area)

$F_W \underline{x_W} + F_M x_M \leq F$ (total amount of fertilizer)

$P_W x_W + P_M x_M \leq P$ (of pesticide)

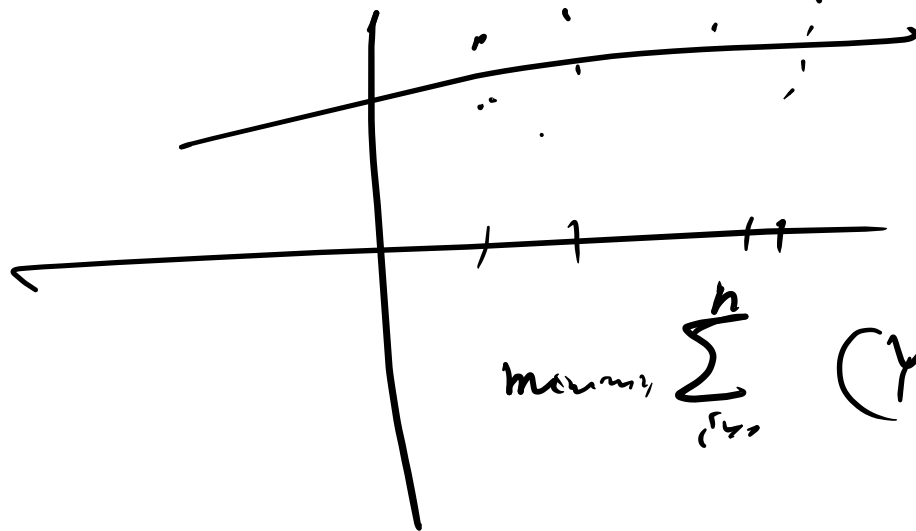
$x_W \geq 0, \quad x_M \geq 0.$

Example 2, (Linear regression)

$(X_1, Y_1) \dots (X_n, Y_n)$

$X \rightarrow$ predictor
variable

$Y \rightarrow$ response variable



$$Y = \beta_0 + \beta_1 X + \epsilon$$

Find best fitting line
to the data.

minimize $\sum_{i=1}^n (Y_i - \beta_0 - \beta_1 X_i)^2$

$$\begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix}$$

$$\begin{bmatrix} \underline{p_0 - p}, \underline{x_1} \\ \vdots \\ \underline{p_0 - p}, \underline{x_n} \end{bmatrix}$$

measure of closeness of the above two

vector.
[l^1 -norm:

l^2 -norm.

l^p -norm.

$$\sum_{i=1}^n |x_i|$$

$$\|\vec{x}\|_1 =$$

$$\sum_{i=1}^n |x_i|$$

$$\|\vec{x}\|_2 =$$

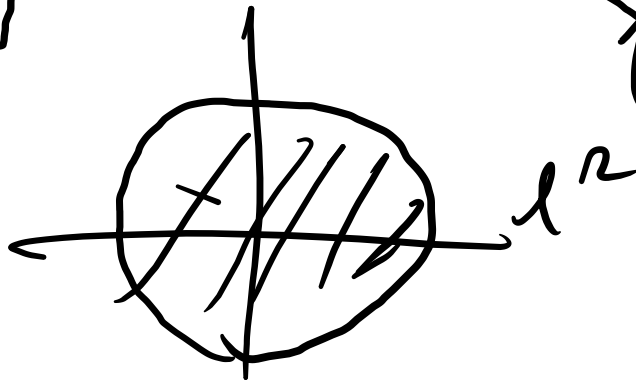
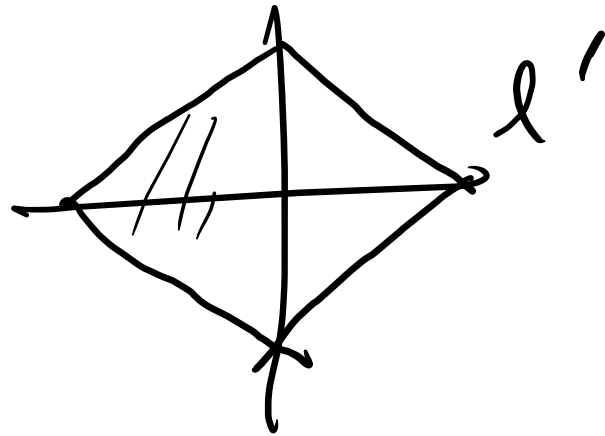
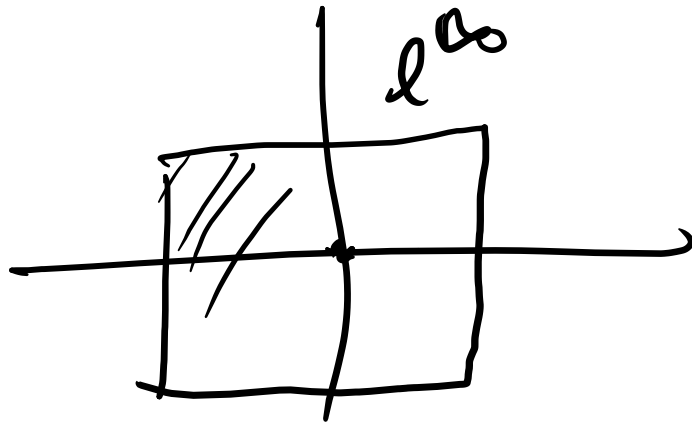
$$\sqrt{\sum_{i=1}^n x_i^2}$$

$$\|\vec{x}\|_p = \left(\sum_{i=1}^n |x_i|^p \right)^{1/p}$$

$$\|\vec{x} - \vec{y}\|_1 = \sum_{i=1}^n |x_i - y_i|$$

(Chebyshev approximation problem)

l^∞ norm $\|x\|_\infty = \max_{1 \leq i \leq n} |x_i|$



$$\text{minimize} \left(\max_{i=1, \dots, n} |Y_i - \underline{\beta}_0 - \underline{\beta}_1 X_i| \right)$$

Equivalent LP:

$$\begin{aligned} &\text{minimize} \quad t \\ &\text{subject to} \quad \underline{\beta}_0 + \underline{\beta}_1 X_i + t \geq Y_i \\ &\quad \quad \quad \underline{\beta}_0 + \underline{\beta}_1 X_i - t \leq Y_i \end{aligned}$$

$$t \geq 0$$

$$\beta_0, \beta_1, t \in \mathbb{R} \\ \geq 0$$

$$\begin{aligned} &i=1, \dots, n \\ &Y_i - \beta_0 - \beta_1 X_i \leq t \\ &[|Y_i - \beta_0 - \beta_1 X_i| \leq t] \end{aligned}$$

GOAL, To understand the structure of
the feasible space in an LP.

$\mathbb{R}^n$

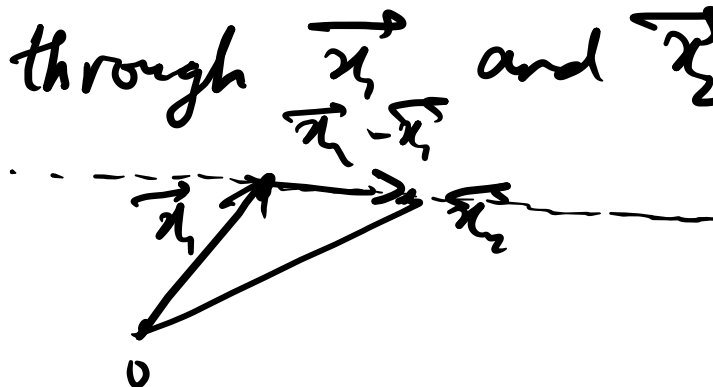
vector subspace $\rightarrow$ closed under arbitrary
linear combinations.

Affine sets in $\mathbb{R}^n$, Convex sets in $\mathbb{R}^n$,
Cones in $\mathbb{R}^n$.

AFFINE SETS

$$\vec{x}_1, \vec{x}_2 \in \mathbb{R}^n \quad (\vec{x}_1 \neq \vec{x}_2)$$

Line passing through $\vec{x}_1$ and $\vec{x}_2$;
given by



$$\vec{x}_1 + t(\vec{x}_2 - \vec{x}_1), \quad t \in \mathbb{R}$$

$$\boxed{t\vec{x}_2 + (1-t)\vec{x}_1}, \quad t \in \mathbb{R}.$$

Line segment joining $\vec{x}_1$ and $\vec{x}_2$:

$$\vec{x}_1 + t(\vec{x}_2 - \vec{x}_1), \quad t \in [0, 1]$$

$$\boxed{t\vec{x}_2 + (1-t)\vec{x}_1}, \quad t \in [0, 1]$$

Affine set: A set $A \subseteq \mathbb{R}^n$ is said to be affine if the unique line through two distinct points in A , lies in A .

In other words,

if $\vec{x}_1, \vec{x}_2 \in A$, then $t\vec{x}_1 + (1-t)\vec{x}_2 \in A$

$\forall t \in \mathbb{R}$.

Equivalently, (prove by induction), ^{linear combination} affine combination of $\vec{x}_1, \dots, \vec{x}_k$

if $\vec{x}_1, \dots, \vec{x}_k \in A$, then $t_1\vec{x}_1 + \dots + t_k\vec{x}_k \in A$

$\forall t_1, \dots, t_k \in \mathbb{R}$ s.t. $\boxed{t_1 + \dots + t_k = 1}$

Remark: If A is an affine set, and $\vec{x}_0 \in A$,

then $\bigcirc \vec{v} : A - \vec{x}_0 = \{ \vec{x} - \vec{x}_0 : \vec{x} \in A \}$
is a vector subspace of $\mathbb{R}^n$ (Exercise).

In other words, every affine subset of $\mathbb{R}^n$ is a translation of a subspace of $\mathbb{R}^n$.

Example: (solution set of a system of linear equations).

$A = \{ \vec{x} : A\vec{x} = \vec{b} \}$, where $A \rightarrow m \times n$ matrix
(m linear equations). $\vec{b} \in \mathbb{R}^m$

$$A\vec{x} = \vec{b}, \quad A\vec{x} = \vec{0}$$

→ particular solution + solution to homogeneous eqⁿ
(translate) (subspace of $\mathbb{R}^n$)

Propⁿ: Every affine set in $\mathbb{R}^n$ can be expressed as the solution set of a system of linear equations.

Proof. A



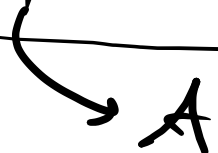
$$V \subset A - \underbrace{x_0}_{\text{circled}}$$

is a subspace of $\mathbb{R}^n$.

* We will express V as the solution set of a system of homogeneous linear equations.

Let $w_1, \dots, w_n$ be an orthonormal basis of V +
(that is, $\langle w_i, w_j \rangle = \delta_{ij}$, $\langle w_i, v \rangle = 0 \quad \forall v \in V$)

$$\boxed{[w_1 \mid w_2 \mid \dots \mid w_k]^T} v = 0$$



 A

Mantra:

Affine set in $\mathbb{R}^n \approx$ translate of a
 subspace of $\mathbb{R}^n \subseteq$ solution set of system
 of linear equations on n variables

Relevance to LP

If the constraint functions defining the decision space were all equalities, the decision space would be an affine subset of $\mathbb{R}^n$.

Affine hull of $S \subseteq \mathbb{R}^n$

Let S be a subset of $\mathbb{R}^n$ (possibly infinite).
The affine hull of S (denoted $\text{aff}(S)$) is the smallest affine set in $\mathbb{R}^n$ containing S , or the

following lemma. If A' is an affine set in $\mathbb{R}^n$ with $S \in A'$ then

$$\text{aff}(S) \subseteq A'$$

$$\left[\text{aff}(S) = \left\{ \theta_1 \vec{x}_1 + \dots + \theta_k \vec{x}_k : \begin{matrix} \vec{x}_1, \dots, \vec{x}_k \in S \\ \theta_1 + \dots + \theta_k = 1 \end{matrix} \right\} \right]$$

↪ set of all affine linear combinations of points in S .

CONVEX SETS

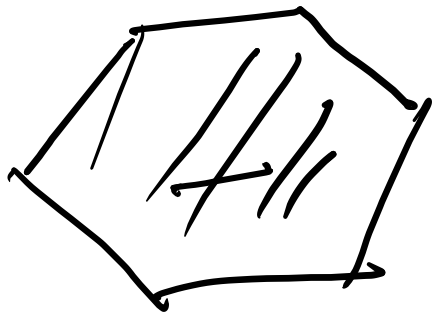
A set $C \subseteq \mathbb{R}^n$ is said to be **convex** if the line segment joining any two distinct points in C , also lies in C .

If $\vec{x}_1, \vec{x}_2 \in C$, then $t\vec{x}_1 + (1-t)\vec{x}_2 \in C$
 $\forall t \in [0, 1]$

Equivalently (prove by induction; exercise)

If $\vec{x}_1, \dots, \vec{x}_n \in C$, then $\theta_1 \vec{x}_1 + \dots + \theta_n \vec{x}_n \in C$
for $\boxed{\theta_i \geq 0}$ and $\boxed{\theta_1 + \dots + \theta_n = 1}$.

* All affine sets in $\mathbb{R}^n$ are convex.



convex



not convex

Convex hull of $S \subseteq \mathbb{R}^n$

Let S be a subset of $\mathbb{R}^n$ (possibly infinite)

The convex hull of S (denoted $\text{conv}(S)$) is the smallest convex set in $\mathbb{R}^n$ containing S , in the following sense;

If C' is any convex set in $\mathbb{R}^n$,
and $S \subseteq C'$, then $\text{conv}(S) \subseteq C'$.

$$\text{conv}(S) = \left\langle \theta_1 \vec{x}_1 + \dots + \theta_k \vec{x}_k \mid \vec{x}_i \in S, \right. \\ \left. k \in \mathbb{N}, \theta_1 + \dots + \theta_k = 1, \theta_i \geq 0 \right\}$$

→ set of all convex combinations of
points in S

Relevance of convex sets - 2P.