# Towards connectivities of random geometric complexes 

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Leiden, June 2018.


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- $\mathcal{P}_{n}=\left\{X_{1}, \ldots, X_{N_{n}}\right\}-X_{i}$ i.i.d. uniform and $N_{n} \sim \operatorname{Poi}(n)$.
- Random geometric graph: $G\left(\mathcal{P}_{n}, r\right)$ : Vertices, $V=\mathcal{P}_{n}$, Edges : $x_{i} \sim x_{j}$ if $0<\left|x_{i}-x_{j}\right| \leq 2 r, r>0$.


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- Threshold for isolated nodes : $J_{n, 0}$. For $w(n) \rightarrow \infty$,

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\mathrm{P}\left(J_{n, 0}=0\right) \rightarrow \begin{cases}0 & \text { if } n \theta_{d} 2^{d} r_{n}^{d}=\log n-w(n) \\ 1 & \text { if } n \theta_{d} 2^{d} r_{n}^{d}=\log n+w(n)\end{cases}
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- $G_{k}^{D}:$ Vertices - $F^{k}$, Edges $\sigma \stackrel{d}{\sim} \sigma^{\prime}$.
- $G_{0}^{U}$ : Vertices - $F^{0} \subset V$. Edges $\subset F^{1}$ i.e., the usual graph.

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C(\mathcal{X}, r) \subset R(\mathcal{X}, r) \subset C(\mathcal{X}, \sqrt{2} r)
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- 1-faces are down-connected but have 2 up-connected components - [1, 2], $\{[2,3],[3,4],[2,4],[4,5],[2,5]\}$.

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- up-connectivity $\nLeftarrow$ homological connectivity.
- Combinatorial structure matters for up/down connectivity !

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- Topological data analysis: Carlsson, 2014 ; Adler, 2015 ; Bobrowski-Kahle, 2017.

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- $G_{k}^{*}\left(X_{n}(r)\right): *=U / D$ i.e., up / down.
$X_{n}(r):=X\left(\mathcal{P}_{n}, r\right), X=C / R$ i.e., Čech / Vietoris-Rips.

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- Also threshold for coverage - ([P. Hall, 1988])


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- $m_{k}=\min \left\{\left|\cap_{i=0}^{k} B_{x_{i}}(2)\right|:\left|x_{i}-x_{j}\right| \leq 2\right\}$.


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- For fixed " $r \in(\log n-\log \log n, \log n)$ ", no isolated edge but infinitely many appear and disappear very quickly !

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- Rate of $c_{n} \rightarrow m_{k} \Rightarrow 2$ nd order term $=k \log \log n$ or $(k-2) \log \log n$ or $(k-1) \log \log n$.

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- Lower bounds involve second-moments and some more Palm calculus.


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- Key geometric Lemma: If $\mathbf{x}, \mathbf{y} \in \mathcal{X}^{(k+1)}$ are isolated $k$-faces $((k+1)$-cliques) in $R(\mathbf{x} \cup \mathbf{y} \cup \mathcal{P}, 1)$ then for some $\epsilon, \beta>0$, either $|Q(\mathbf{x}) \vee Q(\mathbf{y})| \geq m_{k}+\epsilon$ or $|Q(\mathbf{x}) \backslash Q(\mathbf{y})| \geq \beta$.

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- Equivalent to show $\lambda_{2}\left(G\left(\mathcal{P}_{n} \cap B_{1}(O), 1\right)\right)>1 / 2$.


## What more?

- Normalized Laplacian: $L=I-D^{-1 / 2} A D^{-1 / 2}$
- Garland's method: Garland '73, Ballman-Świątkowski '97. : If maximal cliques in a graph $G$ are atleast of order $(k+2)$ and $\lambda_{2}\left(k_{\sigma}\right)>1-1 / k$ for all $k$-cliques $\sigma$, then $\beta_{k-1}(R(G))=0$.
- Used by Kahle '14 to show threshold for $R(G(n, p))$ -Erdös-Rényi clique complex.
- In our case for $k=1$ : To show that w.h.p., for every $X \in \mathcal{P}_{n}$, $\lambda_{2}\left(G\left(\mathcal{P}_{n} \cap B_{2 r_{n}}(X), 2 r_{n}\right)\right)>1 / 2$.
- Equivalent to show $\lambda_{2}\left(G\left(\mathcal{P}_{n} \cap B_{1}(O), 1\right)\right)>1 / 2$.
- "Bobrowski-D.Y." $\lambda_{2}\left(G\left(\mathcal{P}_{n} \cap B_{1}(O), 1\right)\right) \rightarrow 1 / 2$ a.s. !!!


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