

Recall :-

- $\mathcal{S}_N = \{ -1, 1 \}^N$
- $\mathcal{F} = \mathcal{P}(\mathcal{S}_N)$
- $\text{TP}: \mathcal{F} \rightarrow [0, 1]$
- $\text{TP}(A) = \frac{|A|}{2^N} \quad A \subseteq \mathcal{S}_N$
- $\kappa \leq N \quad X_k: \mathcal{S}_N \rightarrow \{-1, 1\}$
- $X_k(\omega) = \omega_k$
- $S_n = \sum_{k=1}^n X_k(\omega), \quad S_0 = n \geq 1 \in \mathbb{N}$
- $\underbrace{\quad}_{\text{finite length random walk}}$
- $P(S_n=x) = P(S_n=-x) = \frac{n!}{\frac{n-x}{2}! \frac{n+x}{2}!} \quad x \in \{-n, -n+2, \dots, n\}$

- Markov chain; independent increments; conditional law
- $A \subseteq \mathcal{S}_n$  - observable by time  $n$  if it was a union of basic events  
 $\{ \omega \in \mathcal{S}_n \mid \omega_1 = \omega_1, \dots, \omega_n = \omega_n \} \quad \omega_i \in \{-1, 1\}$

$\mathcal{A}_n = \{ A \subseteq \mathcal{S}_n \mid A \text{ is observable by time } n \}$   $\phi \in \mathcal{A}_n$   
 $\rightarrow$  closed under complements, unions  $\subseteq$  intersection

Filtration

$$\{\phi, \mathcal{S}\} := \mathcal{A}_0 \subseteq \mathcal{A}_1 \subseteq \dots \subseteq \mathcal{A}_n \subseteq \dots \quad \mathcal{A}_n = \mathcal{F}_n$$

- $T: \mathcal{S} \rightarrow \{0, 1, \dots, N\} \cup \{\infty\}$  is called a stopping time if  
 $\{T=k\} \in \mathcal{A}_k \quad \forall k=0, \dots, N$

Theorem 1.8 [No Profit at favorite times]

Let  $T: \mathcal{S} \rightarrow \{0, 1, \dots, N\} \cup \{\infty\}$  be a stopping time. Then

$$E[S_T] = 0$$

- [Suppose  $\tau$  in notation, i.e.  $\mathcal{S} = \mathcal{S}_0 \cup \dots$  and so on.]

Proof:-

$$S: \Omega \rightarrow \mathbb{Z}$$

$$T: \Omega \rightarrow \{0, 1, \dots, n\}$$

$$S_T: \Omega \rightarrow \mathbb{Z}$$

$$S_{T(\omega)}^{(\omega)} =: S_T^{(\omega)}$$

$$S_0 = 0$$

$$E[S_T] = E\left[\sum_{k=1}^n S_k \mathbf{1}_{(T=k)}\right], \text{ where } \mathbf{1}_A^{(\omega)} = \begin{cases} 1 & \omega \in A \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{aligned} \text{Now, } \sum_{k=1}^n S_k \mathbf{1}_{(T=k)} &= \sum_{k=1}^n S_k [\mathbf{1}_{T \geq k} - \mathbf{1}_{T \geq k+1}] \\ &= \sum_{k=1}^n (S_k - S_{k-1}) \mathbf{1}_{T \geq k} \\ &= \sum_{k=1}^n X_k \mathbf{1}_{T \geq k} \end{aligned}$$

$$\therefore E[S_T] = \sum_{k=1}^n E[X_k \mathbf{1}_{T \geq k}]$$

$1 \leq k \leq n , \quad X_k \mathbf{1}_{T \geq k} = \begin{cases} 1 & X_k = 1 \in T \geq k \\ -1 & X_k = -1 \in T \geq k \\ 0 & \text{otherwise} \end{cases}$

$$\therefore E[X_k \mathbf{1}_{T \geq k}] = P(X_k = 1, T \geq k) - P(X_k = -1, T \geq k)$$

$$\cdot \{T \geq k\}^c = \bigcup_{\ell=1}^{k-1} \{T = \ell\} \in \mathcal{A}_{k-1} \Rightarrow \{T \geq k\}^c \in \mathcal{A}_{k-1}$$

"Stopping time is greater than or equal to k  
should be determined by  
 $X_1, \dots, X_{k-1}$ "

$\cdot X_k$  is independent of  $X_1, \dots, X_{k-1}$

$$\Rightarrow X_k = 1 \text{ is } \perp \text{ of } T \geq k$$

$$X_k = -1 \text{ is } \perp \text{ of } T \geq k$$

$$\{ \begin{aligned} P(X_k=1, T \geq k) &= P(X_k=1) P(T \geq k) \\ P(X_k=-1, T \geq k) &= P(X_k=-1) P(T \geq k) \end{aligned} \} = \frac{1}{2} P(T \geq k)$$

$$1 \leq k \leq \infty$$

$$\begin{aligned} E[X_k \mathbb{1}_{T \geq k}] &= P(X_k=1, T \geq k) - P(X_k=-1, T \geq k) \\ &= \frac{1}{2} P(T \geq k) - \frac{1}{2} P(T \geq k) = 0 \end{aligned}$$

$$\Rightarrow E[S_T] = \sum_{k=1}^{\infty} E[X_k \mathbb{1}_{T \geq k}] = 0$$

□

Question: What is  $\text{Var}[S_T]$ ?  $T$  - stopping time.

Definition 1.8.1: A game System is a sequence of  $\mathbb{R}$ -valued random variables  $\{V_k\}_{k=1}^N$ .

•  $V_k : \Omega \rightarrow \mathbb{R}$  such that

$\{V_k = c\} \in \underline{\mathcal{A}_{k-1}}$   $\forall c \in \mathbb{R}, k=1, \dots, N$

" $V_k$  = amount we bet at start of  $k^{th}$  round".

"Result":  $V_k X_k$   
at  $k^{th}$  round

$\begin{cases} V_k > 0 \\ V_k = 0 \\ V_k < 0 \end{cases}$

(Total gain of the game)  $\Rightarrow S_N = \sum_{k=1}^N V_k X_k$ .

Theorem 1.9 :- [ Profitability is impossible! ]

For any game system  $v_1, v_2, \dots, v_n$  the expected total gain of the game vanishes.

$$E[S_n^v] = 0$$

Proof:-

$$E[S_n^v] = E\left[\sum_{k=1}^n v_k X_{ik}\right]$$

$$= \sum_{k=1}^n E[v_k X_{ik}]$$

$$\text{Range of } v_k \equiv \{c_i^k : 1 \leq i \leq m_k\} \quad \Leftarrow \quad v_k = \sum_{i=1}^{m_k} c_i^k \mathbb{1}(v=c_i^k)$$

$$= \sum_{k=1}^n \sum_{i=1}^{m_k} c_i^k \underbrace{E[X_k \mathbb{1}(v_k=c_i^k)]}_{X_k \perp A_{k-1} \text{ independent}}$$

"as in proof 1"

$$E[S_T] =$$

at  $\oplus$

$$= \sum_{k=1}^n \sum_{i=1}^{m_k} c_i^k \cdot 0$$

$$= 0$$

□

Theorem 1.10

$$T: \Omega \rightarrow \{0, 1, \dots, n\}$$

$$\text{Var}[S_T] = E[T]$$

Proof :-

$$\text{Var}[S_T] = E[S_T^2] - (E[S_T])^2$$

$$E[S_T^2] - 0 \quad \xleftarrow{\text{Theorem 1.8}}$$

$$S_T^2 = \sum_{k=1}^T S_k^2 \mathbb{1}_{\{T=k\}} = \sum_{k=1}^T (S_k^2 - S_{k-1}^2) \mathbb{1}_{\{T \geq k\}} \quad \text{--- } \textcolor{magenta}{\oplus\oplus}$$

Now,

$$\begin{aligned} S_k^2 &= (S_{k-1} + X_k)^2 = S_{k-1}^2 + X_k^2 + 2X_k S_{k-1} \\ &= S_{k-1}^2 + 1 + 2X_k S_{k-1} \end{aligned}$$

$$\Rightarrow S_k^2 - S_{k-1}^2 = 1 + 2X_k S_{k-1} \quad \text{--- } \textcolor{magenta}{\star}$$

$\therefore \textcolor{magenta}{\star}$  into  $\textcolor{magenta}{\oplus\oplus}$  we have

$$\begin{aligned} S_T^2 &= \sum_{k=1}^T (S_k^2 - S_{k-1}^2) \mathbb{1}_{\{T \geq k\}} \\ &= \sum_{k=1}^T (1 + 2X_k S_{k-1}) \mathbb{1}_{T \geq k} \\ &= \sum_{k=1}^T \mathbb{1}_{T \geq k} + 2 \sum_{k=1}^T \underbrace{X_k S_{k-1}}_{V_k} \mathbb{1}_{T \geq k} \end{aligned}$$

$$\therefore S_T^2 = T + 2 \sum_{k=1}^T X_k V_k$$

where  $V_k = S_{k-1} \mathbb{1}_{T \geq k}$ .

$$\therefore E[S_T^2] = E[T] + 2E[S_N^{\vee}]$$

where  $V_k = S_{k-1} \underbrace{1}_{\text{determined}}_{T \geq k}$

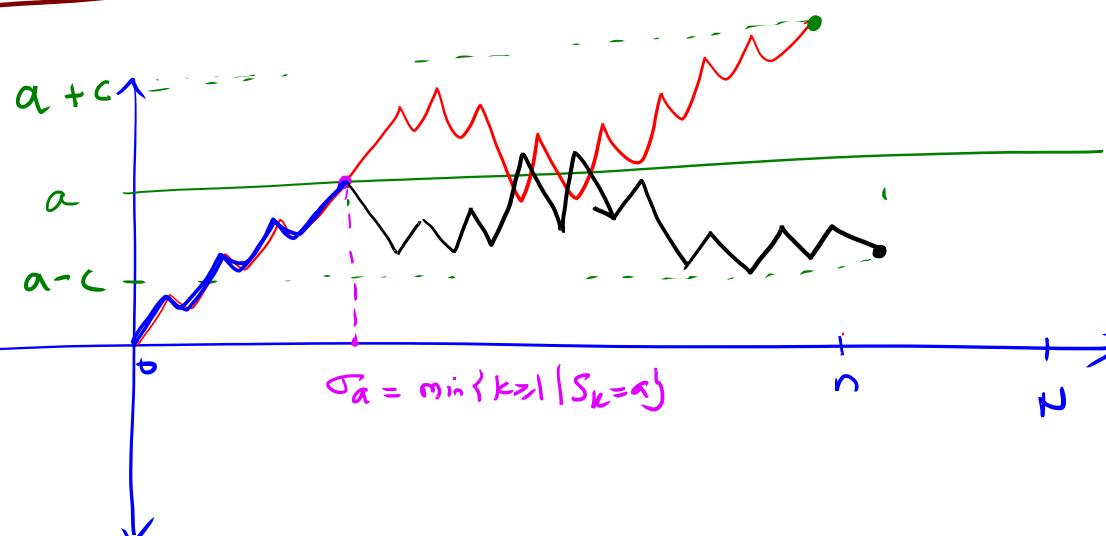
$V_k = c \in A_{k-1}$

$$\therefore \text{Theorem 1.9} \Rightarrow E[S_T^2] = E[T]$$

□

### 1.3 Reflection Principle :-

each path after  
 $\sigma_a$   
 for every real paths  
 $a$  or  $a+c$   
 "reflection"  
 then  $\rightarrow$  Black path  
 $\rightarrow$  at  $a-c$



'Black'  
 III reflection across  
 $y=a$   
 of Red

$c > 0$

$$\begin{aligned} &= |\{S_n = a+c\}| = |\{ \sigma_a \leq n, S_n = a+c \}| \\ &\quad \xrightarrow{\text{has to happen as } S_n = a+c} \\ &= |\{ \sigma_a \leq n, S_n = a-c \}| \end{aligned}$$

Reflection Principle

Lemma 1.11 :-  $a, c \in \mathbb{N}$

$$\mathbb{P}(S_n = a-c, \sigma_a \leq n) = \mathbb{P}(S_n = a+c)$$

Theorem 1.12 : ①  $\mathbb{P}(\sigma_a \leq n) = \mathbb{P}(S_n \notin [-a, a-1])$

(How) ②  $\mathbb{P}(\sigma_a = n) = \frac{1}{2} [\mathbb{P}(S_{n-1} = a-1) - \mathbb{P}(S_{n-1} = a+1)]$

Proof :- ①

$$\mathbb{P}(\sigma_a \leq n) = \sum_{b \in \mathbb{Z}} \mathbb{P}(\sigma_a \leq n, S_n = b)$$

$$= \sum_{\substack{b \in \mathbb{Z} \\ b \geq a}} \mathbb{P}(\sigma_a \leq n, S_n = b) + \sum_{\substack{b \in \mathbb{Z} \\ b < a}} \mathbb{P}(\sigma_a \leq n, S_n = b)$$

"  $b = a-c$  " for  $c > 0$

Reflection

Lemma 1.11  
with  $b = a-c$

$$\begin{aligned} &= \sum_{\substack{b \in \mathbb{Z} \\ b \geq a}} \mathbb{P}(S_n = b) + \sum_{\substack{b \in \mathbb{Z} \\ b < a}} \mathbb{P}(S_n = 2a-b) \\ &= \mathbb{P}(S_n \geq a) + \mathbb{P}(S_n > a) \end{aligned}$$

Symmetry

$$\begin{aligned} &\leftarrow= \mathbb{P}(S_n \geq a) + \mathbb{P}(S_n < -a) \\ &= \mathbb{P}(S_n \notin [-a, a-1]) \end{aligned}$$

□

## 1.4 Arc Sinc law :

$$L = \max \left\{ 0 \leq n \leq 2N : S_n = 0 \right\}$$

"last visit to the origin"  
- NOT a stopping time

Theorem 1.17 :  $n \in \mathbb{N} \setminus \{0\} \quad n \leq N$

$$\mathbb{P}(L = 2n) = \mathbb{P}(S_{2n} = 0) \mathbb{P}(S_{2N-2n} = 0) = 2^{-2n} \binom{2n}{n} \binom{2N-2n}{N-n}$$

Proof - next class.

## Arc-Sinc law :

$$\mathbb{P}(L = 2n) \approx \frac{1}{\pi \sqrt{n(n-r)}}$$

Stirling's formula

$$\mathbb{P}\left(\frac{L}{2N} \leq x\right) = \sum_{n=0, \frac{N}{2} \leq n}^{\infty} \frac{1}{\pi \sqrt{n(n-r)}}$$

...  $n \rightarrow \infty$  ? [Next class]