

Hypothesis Tests and Threshold Crossing Probabilities in Random Walks.

- The MAP Test and the Sequential Test
- Main objective: Bounds on probability
of error in each test
- Tools from Random Walks on \mathbb{Z}
of which is the Chernoff Bound.

- i) We have a discrete signal H taking values $\{0, 1\}$, transmitted over a communication channel. We receive noisy observations of H . Based on data we receive, we want to decide what H is.
- ii) Random Variable $Y \sim P(\cdot | H)$. We want to test for what value H actually is. $H \in \{0, 1\}$.
Based on observations of Y , we calculate a certain statistic

Maximum Apriori Probability Test

Let H be a discrete random variable with support $\{0, 1\}$. We assign

$$P_0 = P(H=0), P_1 = P(H=1).$$

$$\Rightarrow P_0 + P_1 = 1.$$

→ We want to perform a test on H where $H_0: H=0$ vs. $H_1: H=1$.

To perform this test, we make observations y_1, y_2, \dots, y_n of H such that y_1, y_2, \dots, y_n are iid conditional on $H=0$ and iid conditional on $H=1$.

Define $\vec{y} = (y_1, y_2, \dots, y_n)$ and let $f(\cdot | \cdot)$ be the conditional density of y given H .

$$\Rightarrow f(\vec{y} | h) = \prod_{i=1}^n f(y_i | h) \text{ where } h \in \{0, 1\}.$$

From Bayes' Theorem, we know that

$$P(H=0 | \vec{y}) = \frac{p_0 f(\vec{y} | 0)}{\sum_{i=0}^1 p_i f(\vec{y} | i)}.$$

$$P(H=1 | \vec{y}) = \frac{p_1 f(\vec{y} | 1)}{\sum_{i=0}^1 p_i f(\vec{y} | i)}$$

$$\frac{P(H=0 | \vec{y})}{P(H=1 | \vec{y})} = \frac{p_0 f(\vec{y} | 0)}{p_1 f(\vec{y} | 1)}.$$

$$= \frac{P_0}{P_1} \prod_{i=1}^n \frac{f(y_i | 0)}{f(y_i | 1)} = \frac{P_0}{P_1} \lambda(\vec{y})$$

↓
Likelihood
ratio.

The MAP test states that

if $\lambda(\vec{y}) > \frac{P_1}{P_0}$, choose $H=0$

if $\lambda(\vec{y}) \leq \frac{P_1}{P_0}$, choose $H=1$

We want to find bounds on the probability of error in the MAP test.

Why is it important to find a bound on probability of error?

→ Gives us an estimate of how useful our test is.

Define $Z_i = \ln \frac{f(y_i | 0)}{f(y_i | 1)}$ and we

let $\gamma_i(r) = \ln \underbrace{g_{Z_i | H=1}(r)}_{\text{mgf of } Z_i \text{ conditional}} \text{ on } H=1.$

Theorem 1: Let $a = \frac{1}{n} \ln \frac{p_1}{p_0}$ and let $\gamma(r)$ such that $\gamma'(r) = a$.

Then

$$\text{i)} P(\text{choosing } H=0 | H=1) = P(\text{error } | H=1) \\ \leq e^{-n[\gamma(r) - r, a]}$$

$$\text{ii)} P(\text{choosing } H=1 | H=0) = P(\text{error } | H=0) \\ \leq e^{-n[\gamma(r) + (1-r)a]}.$$

Outline of Proof for Theorem 1:

Chernoff Bound: for any random walk $S_n = X_1 + X_2 + \dots + X_n$ where $\{X_i\}_{i \geq 1}$ are iid, let $\gamma_x(r) = \ln g_x(r)$. If γ is such that $\gamma'(s) = a$, then

$$P(S_n \geq na) \leq e^{-n[\gamma(s) - sa]}.$$

This is a sharp bound.

$$\text{i)} P(\text{choosing } H=0 | H=1) \\ = P(\lambda(\vec{y}) > \frac{p_1}{p_0} | H=1) \quad \text{--- } \textcircled{*}$$

$$\text{We let } S_n = \ln(\lambda(\vec{y}))$$

$$\Rightarrow S_n = \sum_{i=1}^n \ln \frac{f(y_i | 1)}{f(y_i | 0)} = \sum z_i$$

S_n is a random walk because

y_1, y_2, \dots, y_n are iid.

This is stated as Proposition 1 in the notes.

$$\textcircled{1} = P(S_n > n\alpha | H=1) \text{ where } \alpha = \frac{1}{n} \ln \frac{p_1}{p_0}$$

From the Chernoff Bound, we have

directly that

$$P(S_n > n\alpha | H=1) \leq e^{n[\lambda(r_1) - r_1 \alpha]}.$$

This proves (i).

For (ii)

$$\begin{aligned} & P(\text{choosing } H=1 | H=0) \\ &= P(\lambda(\vec{y}) \leq \ln \frac{p_1}{p_0} | H=0) \\ &= P(S_n \leq \ln \frac{p_1}{p_0} | H=0) \\ &= P(-S_n \geq -\ln \frac{p_1}{p_0} | H=0) \quad \textcircled{2} \end{aligned}$$

We define $T_n \equiv -S_n$.

$$\Rightarrow T_n = \underbrace{\sum_{i=1}^n x_i}_{s.t. \quad x_i = -z_i}$$

(***)

$$= P(T_n \geq -na \mid H=0)$$

We use the Chernoff bound to see

$$\text{that } P(T_n \geq -na \mid H=0)$$

$$\leq e^{n[\gamma_0(r_0) - r_0(-a)]}$$

$$= e^{n[\gamma_0(r_0) + r_0 n]} \quad (***)$$

$$\text{where } r_0 = \ln g_{X \mid H=0}(r)$$

and r_0 is such that $\gamma_0'(r_0) = -a$.

Claim (proved in the notes)

i) $\gamma_0(r) = \gamma_1(1-r)$ for r that lies
in domain of $\gamma_1(r)$.

ii) $r_0 + r_1 = 1$

iii) $\gamma_0(r_0) = \gamma_1(r_1)$.

We substitute (i), (ii) and (iii)
in (***)) to get

$$P(T_n \geq -na \mid H=0) \\ \leq e^{n[\gamma_1(r_1) + (1-r_1)a]}.$$

This proves (ii)

□.

Chernoff Bound [Proposition 2]

For any random walk $S_n = X_1 + X_2 + \dots + X_n$ where X_i are iid, let $\gamma_x(r) = \ln g_x(r)$. If there exists s such that $\gamma_x'(s) = a$, then $P(S_n \geq na) \leq e^{n[\gamma_x(s) - sa]}$. and this is a sharp bound.

Outline of the proof of Chernoff Bound

$$P(S_n \geq na) = P(e^{rS_n} \geq e^{rna}) \quad \text{for } r > 0.$$

$$\begin{aligned} & P(e^{rS_n} \geq e^{rna}) \\ & \leq \frac{E[e^{rS_n}]}{e^{rna}} \quad [\text{Markov Inequality}] \\ & = [g_x(r)]^n e^{-rna}. \\ \Rightarrow P(S_n \geq na) & \leq [g_x(r)]^n e^{-rna} \\ & = e^{n[\gamma_x(r) - ra]} \quad \text{for } r > 0. \end{aligned}$$

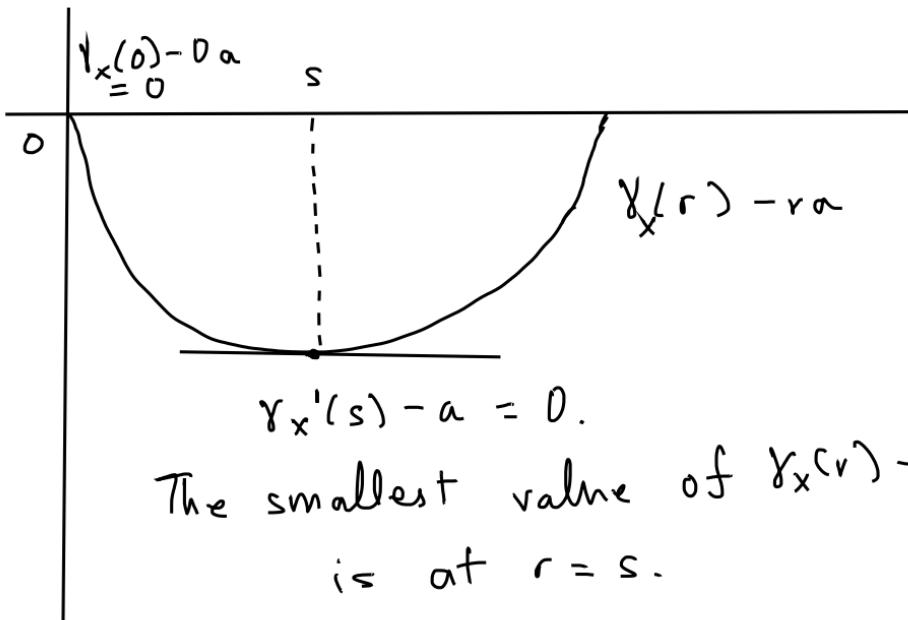
$$\min \gamma_x(r) - ra$$

$$r > 0$$

In theorem 3 of the notes, we use Hölder's Inequality to prove that $\gamma_x(r)$ is a convex function.

If $\gamma_x(r)$ is convex, then a local minimum is also a global minimum.

$$\gamma_x'(s) = a.$$



$$\gamma_x'(s) = a = 0.$$

The smallest value of $\gamma_x(r) - ra$ is at $r = s$.

$$\Rightarrow P(S_n \geq n\alpha) \leq e^{-n[\gamma_{\alpha}(s) - s\alpha]}$$

is the going to be the smallest bound on $P(S_n \geq n\alpha)$. \square .

Sequential test

→ No fixed n .

→ $\alpha > 0, \beta < 0$.

T is a stopping time such that T is the smallest n for which $S_n \geq \alpha$ or $S_n \leq \beta$.

For $n \in \mathbb{N}$, the test states that if

i) If $S_n \in (\beta, \alpha)$, then record observation Y_{n+1}

ii) If $S_n \geq \alpha$ i.e $S_T \geq \alpha$, then we stop testing and we select $H=0$

iii) If $S_n \leq \beta$ i.e $S_T \leq \beta$, then stop testing and we select $H=1$.

Theorem 2:

$$\text{i) } P(\text{choosing } H=0 \mid H=1) = P(\text{error} \mid H=1) \\ \leq e^{-\alpha}$$

$$\text{ii) } P(\text{choosing } H=1 \mid H=0) = P(\text{error} \mid H=0) \\ \leq e^{-\beta}$$

Outline of Proof of Theorem 2

$$\text{i) } P(\text{error} \mid H=1) = P(S_J \geq \alpha \mid H=1) \\ \leq e^{-r^* \alpha} \quad \text{where}$$

$$Y_1(r^*) = 0.$$

Recall that $Y_1(r) = \ln g_{Z \mid H=1}(r)$.

In the notes I have shown that

$$Y_1(1) = 0. \Rightarrow r^* = 1$$

$$\Rightarrow P(S_J \geq \alpha \mid H=1) \leq e^{-\alpha}.$$

This proves (i).

$$\text{ii) } P(\text{error} \mid H=0) = P(S_J \leq \beta \mid H=1) \\ \leq e^{r^* \beta} = e^{\beta}$$

□.

In section 4, I have discussed the power of the MAP test and why the sequential test is advantageous compared to MAP test.

