

Expected hitting time of a pattern from a finite alphabet set

Setup:

let Ω be a finite alphabet set

consider a pattern ψ such that $\psi \in \Omega^m$, for some $m \in \mathbb{N}$

$\{x_n\}_{n \geq 0}$ for $x_n \stackrel{iid}{\sim} \text{Uniform}(\Omega)$

$$T = \min\{k \geq m, k \in \mathbb{N} : (x_{k-m+1}, \dots, x_k) = \psi\}$$

Objective: $E(T)$?

- Example 1: Markov chain techniques, $\psi = \text{ABA}$
- Example 2: Martingale theory, $\psi = \text{ABRACADABRA}$
- Result for arbitrary patterns

Markov chain technique:

Example 1:

$\Omega = \{A, B, C, \dots, X, Y, Z\}$ and $\Psi = ABA$
 X_n is iid Uniform(Ω) and $T = \min\{k \geq 3 : (X_{k2}, X_{k-1}, X_k) = ABA\}$

$\dots \underbrace{***ABA}_{\text{***} \rightarrow **A \rightarrow *AB \rightarrow ABA}$

We partition Ω^3 into $S = \{***, **A, *AB, ABA\}$

let $Y_n = (X_n, X_{n+1}, X_{n+2})$, for $n \geq 0$. We will consider

$\{Y_n\}_{n \geq 0}$ as a Markov chain on S with ABA being the absorbing state

	***	**A	*AB	ABA
***	$\frac{25}{26}$	$\frac{1}{26}$	0	0
**A	$\frac{24}{26}$	$\frac{1}{26}$	$\frac{1}{26}$	0
*AB	$\frac{25}{26}$	0	0	$\frac{1}{26}$
ABA	0	0	0	1

$$|ABA| = 1$$

$$|*AB| = 26$$

$$|**A| = 26^2 - 1$$

$$|***| = 26^3 - 26^2 - 26$$

$$\text{let } T' = \min\{k \geq 0 : Y_k = ABA\} \text{ and}$$

$$h(c) = E(T' | Y_0 = c), \text{ for } c \in \{***, **A, *AB, ABA\}$$

Using one step probability calculation,

References: HW2 - Pb 2

$$h(ABA) = 0$$

$$h(***) = 1 + \frac{25}{26} h(***) + \frac{1}{26} h(**A) + 0 \times h(*AB) + 0 \times h(ABA)$$

$$h(**A) = 1 + \frac{24}{26} h(***) + \frac{1}{26} h(**A) + \frac{1}{26} h(*AB) + 0 \times h(ABA)$$

$$h(*AB) = 1 + \frac{25}{26} h(***) + 0 \times h(**A) + 0 \times h(*AB) + \frac{1}{26} h(ABA)$$

$$h(***) = 26^3 + 26$$

$$h(**A) = 26^3$$

$$h(*AB) = 26^3 - 26^2 + 26$$

$$h(ABA) = 0$$

gt takes 3 steps to reach the initial state c

$$E(T) = E(T') + 3$$

$$E(T') = \sum_{c \in S} P(Y_0 = c) E(T' | Y_0 = c) = \sum_{c \in S} P(Y_0 = c) h(c)$$

$$\text{But, } P(Y_0 = c) = \frac{|c|}{|S|} = \frac{|c|}{26^3} \text{ and}$$

$$P(Y_0 = ABA) = \frac{1}{26^3} ; \quad P(Y_0 = *AB) = \frac{26}{26^3}$$

$$P(Y_0 = **A) = \frac{26^2 - 1}{26^3} ; \quad P(Y_0 = ***) = \frac{26^3 - 26^2 - 26}{26^3}$$

$$\text{Then, } E(T') = 0 \times \frac{1}{26^3} + \frac{26(26^3 - 26^2 + 26)}{26^3} + \frac{(26^2 - 1)(26)}{26^3} + \frac{(26^3 - 26^2 - 26)(26^3 + 26)}{26^3}$$

$$= E(T') = 26^3 + 26 - 3$$

$$\therefore E(T) = 26^3 + 26$$

Pros:

- Setup as a Markov chain with ψ being the absorbing state.
- Calculate transition probabilities and category sizes.
- Do one step probability calculation to find $h(c)$.
- Calculate $E(T')$ and then $E(T)$.

Disadvantages:

- For pattern of length n , need to solve $(n+1) \times (n+1)$ system of linear equations.
- Can get very calculation intensive.

Martingale Method:

Example 2:

$$\Omega = \{A, B, C, \dots, X, Y, Z\} \text{ and } \Psi = \text{ABRACADABRA}$$

$$X_n \text{ is iid Uniform}(\Omega) \text{ and } T = \min_{k \geq 1} \{ (X_{k-10}, \dots, X_k) = \Psi \}$$

Consider a casino with a monkey and a typewriter inside it. The typewriter has just the 26 english alphabets on it and the monkey types one letter randomly every second.


$$u_1, u_2, u_3, u_4, u_5, u_6, u_7, u_8, u_9, u_{10}, u_{11}, u_{12}, \dots$$

At every time k , $k \geq 1$, a new gambler walks into the casino and bets $\frac{1}{k}$ that X_k , the next letter the monkey will type is going to be 'A'. If they are correct, they win $\frac{2}{k}$ otherwise they leave the game.

Why 26? Expected profit = $\frac{1}{26}(26-1) + \frac{25}{26}(-1) = 0$

If the gambler was correct, then in the next time step (next round) they bet the entire ₹26 that the upcoming letter, x_{k+1} will be 'B' (for a prize of ₹26²) and so on.

This goes on until the word ABRACADABRA appears and then the game stops.

Example: In the case when $x_1 = A$, $x_2 = B$, $x_3 = A$, $x_4 = D$, $x_5 = B \dots$

[illegible]

let Z_n^k be the amount of money the k^{th} gambler has at the end of time n

$$Z_n^k = \begin{cases} 1 & , \text{ if } n < k \\ 26 Z_{n-1}^k \cdot \mathbb{1}_{\{X_n = \psi_{n-k+1}\}} & , \text{ if } k \leq n \leq k+10 \\ Z_{k+10}^k & , \text{ if } n \geq k+11 \end{cases}$$

Then, the net profit the casino makes from the k^{th} gambler at the end of the n^{th} round is: $Z_0^k - Z_n^k = 1 - Z_n^k$

Therefore, the total profit the casino makes from all the gamblers, at the end of the n^{th} round is:

$$M_n = \sum_{k=1}^n (1 - Z_n^k) = n - \sum_{k=1}^n Z_n^k$$

$\{M_n\}_{n \geq 1}$ is a martingale w.r.t the filtration $\mathcal{F}_n$

defined by the set of observable events till time n

Since $E(|M_n|) < \infty$, $\forall n \in \mathbb{N}$ and $E(M_n | \mathcal{F}_{n-1}) = M_{n-1}$

M_n is a function of $(X_1, X_2, \dots, X_n)$

- T is a stopping time w.r.t $\{M_n\}_{n \geq 1}$

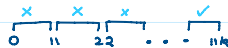
We will find $E(T)$ using $E(M_T)$
How to find $E(M_T)$? We will use OST.

To be able to use OST, we have to show:

- ① $E(T)$ is finite
- ② $|M_n - M_{n-1}| \leq c$, $\forall n \geq 1$, that is, M has bounded increments

To show $E(T)$ is finite:

Consider blocks of size 11



Success if block = 4 and failure otherwise

S is the number of blocks until the first success

So, $S = k \Rightarrow T \leq 11k$



Hence, $T \leq 11S \Rightarrow E(T) \leq 11E(S)$

But $S \sim \text{Geometric}(\frac{1}{26}) \Rightarrow E(T) \leq 11(26)$

$\therefore E(T) < \infty$ ①

To show M has bounded increments:

$$\text{For } n > 1, \text{ consider } |M_n - M_{n-1}| = |n - \sum_{k=1}^n Z_n^k - n + 1 + \sum_{k=1}^{n-1} Z_{n-1}^k|$$

$$\leq \left| \sum_{k=1}^{n-1} (Z_{n-1}^k - Z_n^k) \right| + |Z_n^n - 1|$$

$$\leq \left| \sum_{k=n-10}^{n-1} (Z_{n-1}^k - Z_n^k) \right| + 25 \quad [\text{Since } Z_{n-1}^k = Z_n^k, \text{ for } n \geq k+10 \text{ and } Z_n^n \leq 26]$$

$$= \left| \sum_{k=n-10}^{n-1} (Z_{n-1}^k - 26 Z_{n-1}^k \mathbb{1}_{\{U_n = W_{n-k+1}\}}) \right| + 25$$

$$\leq 25(1 + 26 + 26^2 + \dots + 26^{10}) \quad [\text{Since } Z_n^{n-k+1} \leq 26^k]$$

$$\Rightarrow |M_n - M_{n-1}| \leq 26^{11} - 1, \forall n \in \mathbb{N} \text{ and hence, } M \text{ has bounded increments.} - (2)$$

So, ① and ② $\Rightarrow$ The conditions for OST, i.e., $E(M_T | T > n) P(T > n) \rightarrow 0$ as $n \rightarrow \infty$ and $E(M_T) < \infty$

$$\text{Hence, } E(M_T) = E(M_1)$$

$$= 1 - E(Z_1) = 1 - \left(\frac{1}{26}(26) + \frac{25}{26}(0) \right)$$

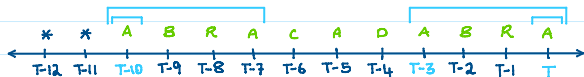
$$\Rightarrow E(M_T) = 0 \quad (3)$$

How are M_T and T related?

$$M_T = T - \sum_{k=1}^T Z_T^k = T - \sum_{k=T-10}^T Z_T^k$$

This is indeed the case, as for all gamblers who started

betting before time $T-10$, $Z_T^k = 0$



Clearly, $Z_T^k = 0$ for $T-10 \leq k \leq T$ and $k \notin \{T, T-3, T-10\}$.

$$\text{Then, } Z_T^T = 26, Z_T^{T-3} = 26^4, Z_T^{T-10} = 26^{11}$$

$$\therefore M_T = T - 26^{11} - 26^4 - 26$$

$$\text{Therefore, } E(M_T) = 0 = E(T - 26^{11} - 26^4 - 26)$$

$$\text{Hence, } E(T) = 26^{11} + 26^4 + 26$$

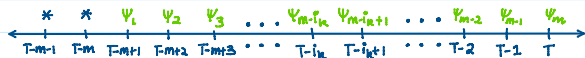
General Result:

x_n i.i.d Uniform(Ω) and $\Psi \in \Omega^m$, for some $m \in \mathbb{N}$.

$T = \min\{k \geq m, k \in \mathbb{N} : (x_{k-m+1}, \dots, x_{k-1}, x_k) = \Psi\}$

Theorem 1: Consider Ψ , a pattern of length $m \in \mathbb{N}$. Suppose there exist $0 < i_1 < i_2 < \dots < i_j < m$, for $0 \leq j < m$, such that, the first i_k letters of the pattern Ψ is equal to the last i_k letters of the pattern Ψ , for $\forall 1 \leq k \leq j$. That is, $(\Psi_1, \Psi_2, \dots, \Psi_{i_k}) = (\Psi_{m-i_k+1}, \Psi_{m-i_k+2}, \dots, \Psi_m)$, $\forall 1 \leq k \leq j$. Then,

$$E(T) = |\Omega|^m + \sum_{k=1}^j |\Omega|^{i_k}$$



Example: For $\Psi = \text{HAHANA}$, $E(T) = 26^6 + 26^4 + 26^2$



But for $\Psi = \text{ABCDEF}$, $E(T) = 26^6$

Sketch of proof:

- Casino with monkey and typewriter setup as before.
- Casino closes when the word Ψ appears and incoming gamblers bet on the letters of the word Ψ .
- Z_n^k defined appropriately and $M_n = n - \sum_{k=1}^j Z_n^k$
- Show that $\{M_n\}_{n \geq 1}$ is a Martingale w.r.t $\mathcal{F}_n$
- Show that $E(T) < \infty$
- Show that M has bounded increments, i.e., $|M_n - M_{n-1}| \leq C$, $\forall n \geq 1$
- Using OST, we have $E(M_T) = E(M_1) = 0$
- Show that $M_T = T - |\Omega|^m - \sum_{k=1}^j |\Omega|^{i_k}$
- And hence, $E(T) = |\Omega|^m + \sum_{k=1}^j |\Omega|^{i_k}$