

Finding zeros of $f: \mathbb{R} \rightarrow \mathbb{R}$ or $f: \mathbb{R}^n \rightarrow \mathbb{R}$

- $f: \mathbb{R} \rightarrow \mathbb{R}$
 - Zeros of polynomials there are known methods
 - if f is not a polynomial then finding a zero(s) becomes non-trivial
- Optimization:- Any minimization / maximization problem reduces to understanding critical points. More precisely

$$F: \mathbb{R}^n \rightarrow \mathbb{R} \quad \max_{x \in A} F(x) \quad \text{or} \quad \min_{x \in B} F(x)$$

One needs to find
 $\{x : \nabla F(x) = 0\}$

we will explore

$$f: \mathbb{R} \rightarrow \mathbb{R}$$

- Bisection method
- Newton-Raphson method
- Secant method
 - Convergence
 - rate
 - limitations
- Both Theoretical aspects & implementation on $\mathbb{R}$

$$F: \mathbb{R}^n \rightarrow \mathbb{R}; \min_{x \in A} F(x)$$

- Golden Section Search
- Generalization of Newton-Raphson
- Gradient descent
- Hill climbing
- Simulated Annealing
- Method, usefulness & implementation on $\mathbb{R}^n$

Bisection Method:-

$f: [a, b] \rightarrow \mathbb{R}$

Assume $f(a) < 0$, $f(b) > 0$, f is continuous.

Aim:- find $c: f(c) = 0$, $c \in (a, b)$

Algorithm:

$a = \dots$

$b = \dots$

for k in $1:n$

$x_k = \frac{a+b}{2}$ [bisect the interval]

if $\text{sign } f(x_k) = \text{sign } f(a)$

$a = x_k$

if $\text{sign } f(x_k) = \text{sign } f(b)$

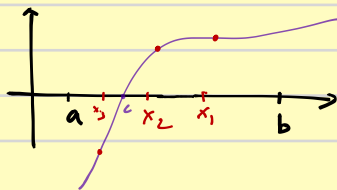
$b = x_k$

end

return (x_n) .

One can decide on n by the following

- stop at some arbitrary large n
- stop when $|f(x_k)| < \text{tol}$



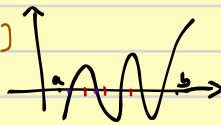
let $c \in (a, b)$ $f(c) = 0$. observe

$$|x_k - c| \leq |a_k - b_k| = \frac{1}{2^k} |a - b|$$

• Each bisection step - discovers a new correct digit in binary expansion of c

• Clearly $x_k \rightarrow c$ as $k \rightarrow \infty$; the method will converge under hypothesis

• [no control]



- could converge to one of the roots

• [NOT really fast] only guarantees that

$$|x_{k+1} - c| < \frac{1}{2} |x_k - c|$$

[fixed]

[linear convergence] will explain below

Extensions:- $f: \mathbb{R} \rightarrow \mathbb{R}$ continuous; one can use bracketing to designate an interval to apply bisection.

Algorithm:- f $x_{\min}, x_{\max}, n$

$$d = \frac{x_{\max} - x_{\min}}{n}$$

$$a = x_{\min}$$

while $i < n$

$$i \rightarrow i + 1$$

$$b = a + d$$

if f changes sign in $[a, b]$, save

end

$$a = b$$

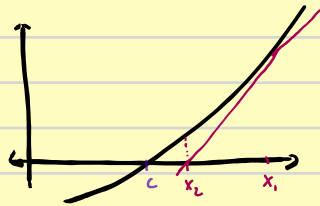
end

Newton-Raphson

- important method and it is the basis of most optimization solutions.

Assume: $f: \mathbb{R} \rightarrow \mathbb{R}$ and f is twice differentiable

Idea:-



- Drop tangent line at $(x_1, f(x_1))$
- x -intercept of tangent line $= x_{n+1}$

Algorithm :-

$x_1 =$ initialize

for k in $2:n$

$$x_k = x_{k-1} - \frac{f(x_{k-1})}{f'(x_{k-1})}$$

at convergence criteria stop

end

return (x_k)

$$y - f(x_1) = (x - x_1) f'(x_1)$$

$(y=0)$ $x_2 = -\frac{f(x_1)}{f'(x_1)} + x_1$

Example (already seen): $f(x) = x^2 - a \quad a > 0$

$$f'(x) = 2x$$

Newton's Method reduces to: $x_1 > \sqrt{a}$

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$

$$\Rightarrow x_{k+1} = x_k - \frac{x_k^2 - a}{2x_k}$$

$$\Rightarrow x_{k+1} = \frac{1}{2} \left(x_k + \frac{a}{x_k} \right)$$

Already seen the following:

• $x_k \rightarrow \sqrt{a} \quad \text{as } k \rightarrow \infty$

$$\& \quad x_{k+1} - \sqrt{a} \leq \frac{(x_k - \sqrt{a})^2}{2\sqrt{a}} \quad \forall k \geq 1$$

[compare with Bisection]
- Quadratic Convergence
- will explain later

• In the case $a=3$, $x_1=2$ we saw

$$x_6 - \sqrt{a} < 10^{-3}$$

Q: - Is this characteristic of the method?

• Assume: $\left\{ \begin{array}{l} f: \mathbb{R} \rightarrow \mathbb{R} \quad ; \quad f \text{ - twice differentiable} \\ - f(c) = 0 \\ - \exists \delta > 0 : f'(x) \neq 0 \quad \forall x \in (c-\delta, c+\delta) \\ - f'' \text{ - continuous.} \end{array} \right.$

If $x_k \in (c-\delta, c+\delta) \quad \forall k \geq 1$

we also have $x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)} \quad (2)$
need $f'(x_k) \neq 0 \rightarrow$

By Taylor's Theorem $\exists \xi_k$ between x_k and c
such that
$$f(c) = f(x_k) + (c-x_k)f'(x_k) + \frac{1}{2}(c-x_k)^2 f''(\xi_k) \quad (1)$$

$$f(c) = 0 \quad (3)$$

Substitute:

③ in ① & value of $f(x_k)$ from ② into ①
to get

$$\begin{aligned} 0 &= (x_{k+1} - x_k)f'(x_k) + (c-x_k)f'(x_k) + \frac{1}{2}(c-x_k)^2 f''(\xi_k) \\ \Rightarrow (x_{k+1} - c) &= \frac{1}{2} \frac{f''(\xi_k)}{f'(x_k)} (x_k - c)^2 \quad (4) \end{aligned}$$

$$\Rightarrow (x_{k+1} - c) \leq M_1 (x_k - c)^2$$

$$\Rightarrow (x_{k+1} - c) \leq \frac{[M_1 (x_0 - c)]^{2^k}}{M_1}$$

If $|M_1 (x_0 - c)| < 1$ then $x_k \rightarrow c$
as $k \rightarrow \infty$.

From the above calculation we can postulate

Theorem (a Sufficient Condition) :-

$f: \mathbb{R} \rightarrow \mathbb{R}$; f - twice differentiable

• $f(c) = 0$

• $\exists 0 < \delta, f'(x) \neq 0 \quad \forall x \in (c-\delta, c+\delta)$

• f'' - continuous in $(c-\delta, c+\delta)$

$\exists \eta > 0$ st. if $x_0 \in (c-\eta, c+\eta)$

(a)
$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)} \quad [\text{well defined } \forall k \geq 1]$$

(b)
$$(x_{k+1} - c) < M (x_k - c)^2$$

$$\exists \quad x_{k+1} - c \leq \underbrace{(M (x_0 - c))}_{M}^{2^k}$$

for any $M > \frac{1}{2} \frac{|f''(c)|}{|f'(c)|}$

In particular $x_k \rightarrow c$ as $k \rightarrow \infty$
(quadratically).

Secant Method ..

$f: \mathbb{R} \rightarrow \mathbb{R}$, Suppose computing the derivatives are hard. We discussed last week, finite difference approximation to derivative.

$x, y \in \mathbb{R}, x \neq y$ set

$$f[x, y] := \frac{f(x) - f(y)}{x - y}$$

Idea:-



- Draw a secant line from x_0, x_1
- x -intercept of line = x_2

$$x_2 = x_1 - \frac{f(x_1)}{f[x_1, x_0]}$$

Algorithm :-

$x_1 =$ initialize

for k in $2:n$

$$x_k = x_{k-1} - \frac{f(x_{k-1})}{f[x_{k-1}, x_{k-2}]}$$

$$f[x_{k-1}, x_{k-2}]$$

at convergence criteria stop

end

return (x_k)

Assume: $\left\{ \begin{array}{l} f: \mathbb{R} \rightarrow \mathbb{R} \text{ ; } f \text{ - twice differentiable} \\ - f(c) = 0 \\ - \exists \delta > 0 : f'(x) \neq 0 \quad \forall x \in (c - \delta, c + \delta) \\ - f'' \text{ - continuous.} \end{array} \right.$

$$n \geq 1 \quad x_{n+1} = x_n - \frac{f(x_n)}{f[x_n, x_{n-1}]} \quad - (1)$$

$$\text{let } p(x) = f(x_n) + f[x_{n-1}, x_n] (x - x_n)$$

$$f(x) - p(x)$$

$$= f(x) - \left[f(x_n) + \frac{f(x_n) - f(x_{n-1})}{x_n - x_{n-1}} (x - x_n) \right]$$

$$= \left[\frac{f(x) - f(x_n)}{x - x_n} - \frac{f(x_n) - f(x_{n-1})}{x_n - x_{n-1}} \right] (x - x_n)$$

$$= f[x, x_n, x_{n-1}] (x - x_n) (x - x_{n-1})$$

$$\stackrel{\text{claim } (*)}{=} \frac{f''(\xi)}{2} (x - x_n) (x - x_{n-1})$$

where $\xi \in [a, b] \supseteq \{x, x_n, x_{n-1}\}$.

Applying above at $x=c$

$$0 = f(x_n) + f[x_{n-1}, x_n] (c - x_n) + \frac{f''(\xi)}{2} (c - x_n) (c - x_{n-1}) \quad - (2)$$

(1) yields

$$0 = f(x_n) + f[x_{n-1}, x_n] (x_{n+1} - x_n) \quad - (3)$$

$$(2) - (3) \Rightarrow$$

$$c - x_{n+1} = \frac{1}{2} \frac{f''(\xi)}{f[x_{n-1}, x_n]} (c - x_n) (c - x_{n-1})$$